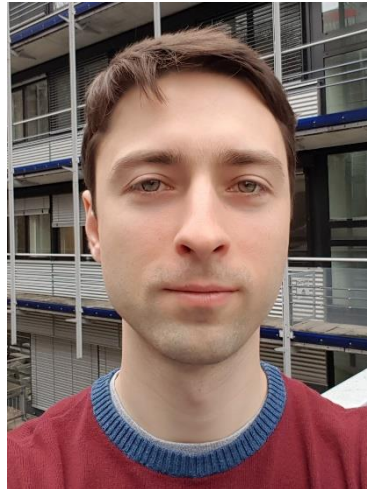


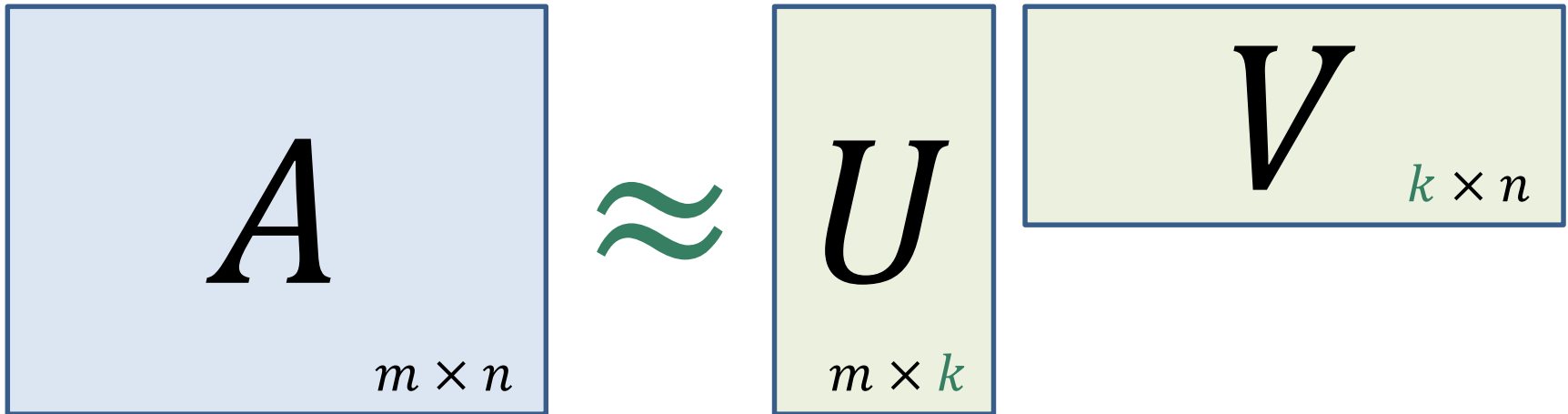
Approximation Algorithms for ℓ_0 -Low Rank Approximation

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Low Rank Approximation



Motivation:

- Compact Storage
- Faster Matrix-Vector Multiplication
- Removes Outliers/Anomalies

} $mn \mapsto k(m + n)$

ℓ_0 -Low Rank Approximation

$$\min_{\substack{U \in \mathcal{S}^{m \times k} \\ V \in \mathcal{S}^{k \times n}}} \left| \begin{array}{c} \boxed{A} \\ \in \mathcal{S}^{m \times n} \end{array} \right. - \left. \begin{array}{c} \boxed{U} \\ m \times k \end{array} \begin{array}{c} \boxed{V} \\ k \times n \end{array} \right| \mathbf{0}$$

Input: semiring/field \mathcal{S} , matrix $A \in \mathcal{S}^{m \times n}$, positive integer k

Output: rank- k matrix decomposition U, V minimizing the
of disagreeing components -- $|M|_0 = \text{nnz}(M)$

ℓ_0 -Low Rank Approximation

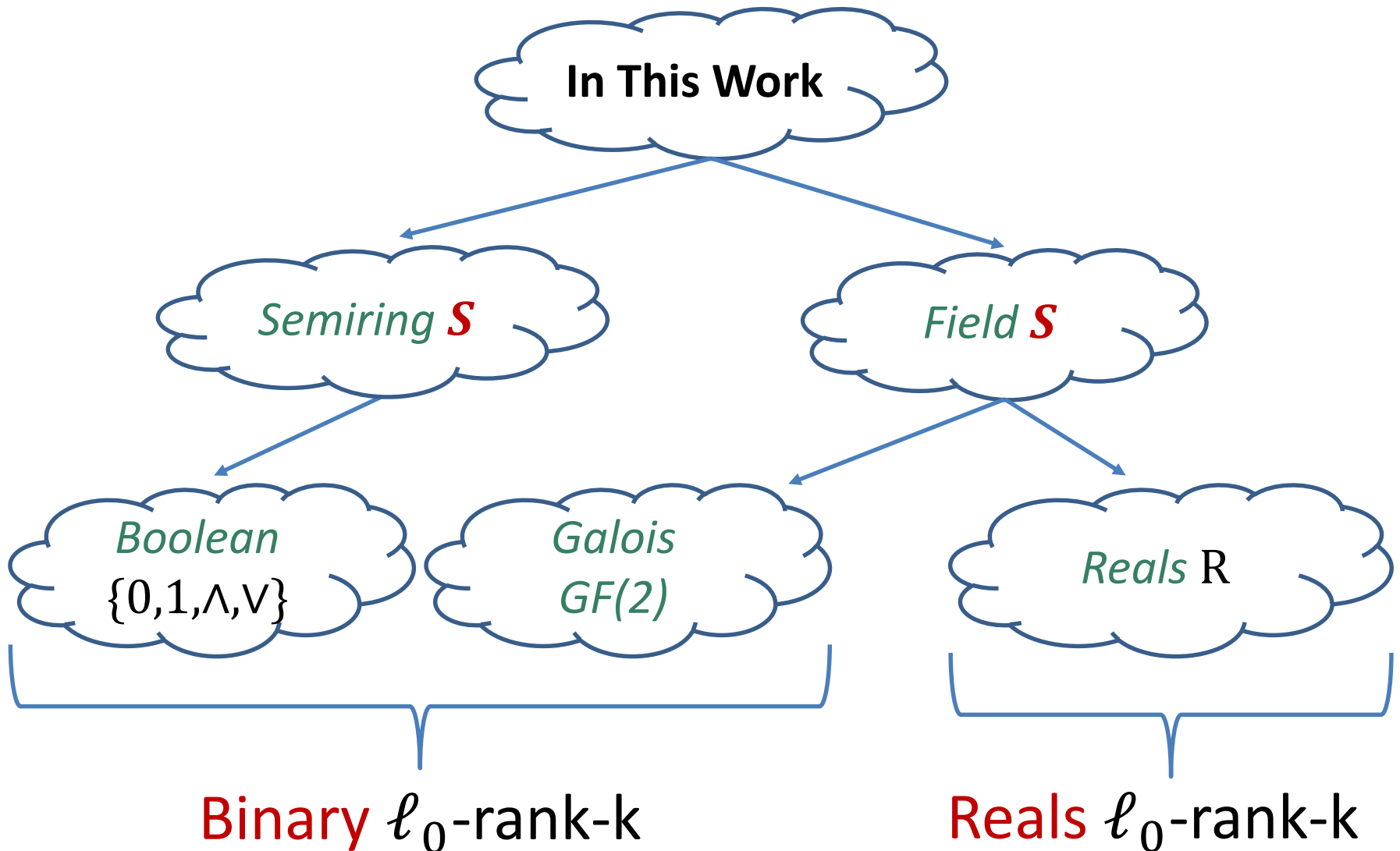
$$\min_{\substack{U \in S^{m \times k} \\ V \in S^{k \times n}}} \left| \begin{array}{c} \boxed{A} \\ \in S^{m \times n} \end{array} - \begin{array}{c} \boxed{U} \\ m \times k \end{array} \begin{array}{c} \boxed{V} \\ k \times n \end{array} \right|_0$$

Why $|A|_0$?

1. *Robust* to outliers/anomalies.

2. Natural for problems with *no underlying metric*.

ℓ_0 -Low Rank Approximation



The ℓ_0 -Low Rank Approximation

Classifications based on S

Our Result:

Poly-time Bicriteria Algorithm

approximately solves
original Robust PCA
by *Candes et al.*

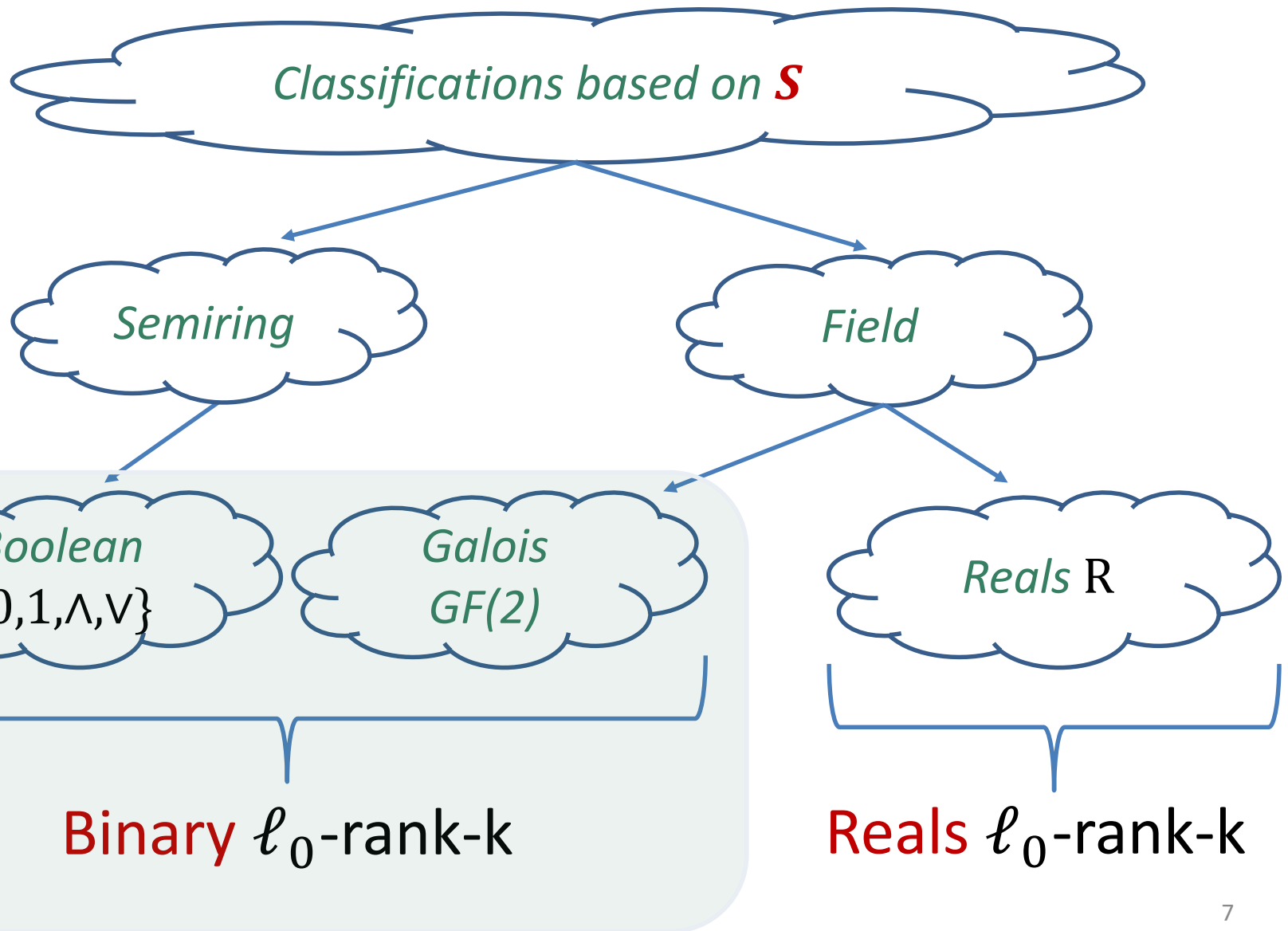
Field

Reals \mathbb{R}

Binary ℓ_0 -rank-k

Reals ℓ_0 -rank-k

The ℓ_0 -Low Rank Approximation



Binary ℓ_0 -rank- k

$$A \in \{0,1\}^{m \times n}$$

$$\operatorname{argmin}_{U \in \{0,1\}^{m \times k}, V \in \{0,1\}^{k \times n}} \|A - U \cdot V\|_0$$

Known under many different names:

Matrix Rigidity

– Complexity Theory

Boolean Factor Analysis

– Machine Learning & Data Mining

Binary Matrix Factor

Biclique (k-way) Partition – Graph Theory

NP-Hard even for $k = 1$

Binary ℓ_0 -rank-1

$$A \in \{0,1\}^{m \times n}$$

Our Algorithmic Results

$$\text{OPT} = \min_{u \in \{0,1\}^m, v \in \{0,1\}^n} |A - u \cdot v|_0 \leq |A|_0$$

Small $\phi = \frac{\text{OPT}}{|A|_0}$

Exact Alg.

Runtime: $O(|A|_0 + m + n)$

Returns: $(1 + \phi)$ – *Approx.*

Good Approx: $A \approx \text{rank-1 matrix}$

Runtime: $2^{\text{OPT}/\sqrt{|A|_0}} \text{poly}(nm)$

Returns: An **Optimal** Solution

PolyTime: $\text{OPT} \leq O(\log(nm))\sqrt{|A|_0}$

Binary ℓ_0 -rank- k

$A \in \{0,1\}^{m \times n}$

$$\operatorname{argmin}_{U \in \{0,1\}^{m \times k}, V \in \{0,1\}^{k \times n}} \|A - U \cdot V\|_0$$

*NP-Hard
even for
 $k = 1$*

*Does the Binary ℓ_0 -rank- k
admits a **Polynomial Time
Approximation Scheme
(PTAS)** ?*

*In a follow up work,
we give an **affirmative** answer.*

Thank You



Carnegie
Mellon
University

